

CSE 462 : Relational Model and Relational Algebra

Name: _____

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***** Solved in class, on January 31, 2011. *****

1. Consider the relation instance R below.

A	B	C
1	2	3
1	3	4
2	3	4
2	3	5

a) Specify a schema for this relation, including a minimal key.

$R(\underline{A}: \text{integer}, \underline{B}: \text{integer}, \underline{C}: \text{integer})$

b) Based on your answer to 1a), explain whether the tuple (2,4,5) can be inserted into R .

It cannot. It would violate the key constraint: no two tuples may agree on their A and C attribute values, but the new tuple and the existing tuple (2,3,5) would.

c) In how many different ways can R be presented?

3! · 4!

2. Consider the relation schema $\text{Employee}(\underline{\text{SSN}}, \text{LastName}, \text{FirstName}, \text{DoB}, \text{Salary}, \text{ManagerSSN})$. Provide a CREATE TABLE statement to create this relation using reasonable data types and constraints to guarantee consistency of the relation.

```
CREATE TABLE Employee (
  SSN          CHAR(9) NOT NULL PRIMARY KEY,
  LastName     VARCHAR(100) NOT NULL,
  FirstName    VARCHAR(100) NOT NULL,
  DoB         DATE NOT NULL,
  Salary       NUMERIC NOT NULL CHECK(Salary >= 0),
  ManagerSSN  CHAR(9) REFERENCES Employee(SSN)
);
```

3. Suppose relations R and S have n tuples and m tuples, respectively. Give the minimum and maximum number of tuples that the results of the following expressions can have. Provide your answer in terms of m and n . (Exercise 2.4.7 in the textbook)

	$R \cup S$	$R \cap S$	$R \bowtie S$	$\sigma_{\theta}(R) \times S^*$	$\pi_L(R) - S^{\dagger}$
min	$\max(n, m)$	0	0	0	0
max	$n + m$	$\min(n, m)$	$n \cdot m$	$n \cdot m$	n

* θ is an arbitrary selection formula. $\dagger L$ is an arbitrary projection list.

4. The following algebraic laws hold for sets but not for bags. For each of the laws, show that they do not hold for bags by providing a counterexample in terms of relation instances R , S , and T . (Based on exercise 5.1.5 in the textbook)

a) $(R \cap S) - T = R \cap (S - T)$

b) $R \cap (S \cup T) = (R \cap S) \cup (R \cap T)$

$R(X) = \{(x)\}$, $S(X) = \{(x), (x)\}$, $T(X) = \{(x)\}$ is a counterexample for both (a) and (b).

$((\{(x)\} \cap \{(x), (x)\}) - \{(x)\}) \neq \{(x)\} \cap ((\{(x), (x)\}) - \{(x)\}))$

$\{(x)\} \cap ((\{(x), (x)\} \cup \{(x)\}) \neq ((\{(x)\} \cap \{(x), (x)\}) \cup ((\{(x)\} \cap \{(x)\}))$

$\{(x)\} - \{(x)\} \neq \{(x)\} \cap \{(x)\}$

$\{(x)\} - \{(x)\} \neq \{(x)\} \cap \{(x)\}$

$\emptyset \neq \{(x)\}$

$\emptyset \neq \{(x)\}$